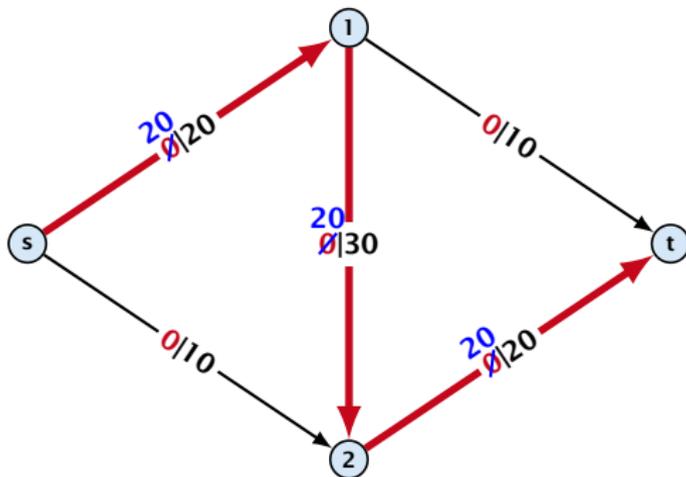# 12

Greedy-algorithm:

- ▶ start with $f(e) = 0$ everywhere
- ▶ find an s-t path with $f(e) < c(e)$ on every edge
- ▶ augment flow along the path
- ▶ repeat as long as possible

The Residual Graph

From the graph $G = (V, E, c)$ and the current flow f we construct an auxiliary graph $G_f = (V, E_f, c_f)$ (the residual graph):

▶ Suppose the original graph has edges $e_1 = (u, v)$, and $e_2 = (v, u)$ between u and v.

▶ G_f has edge e_1' with capacity $\max\{0, c(e_1) - f(e_1) + f(e_2)\}$ and e_2' with with capacity $\max\{0, c(e_2) - f(e_2) + f(e_1)\}$.

G $\quad$ u ← ——— v
$\qquad\qquad$ 14|16 $\qquad$ 10|20

G_f $\quad$ u ← ——— v
$\qquad\qquad$ 12 $\qquad$ 24
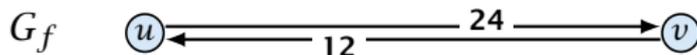
Augmenting Path Algorithm

Definition 1

An augmenting path with respect to flow f, is a path from s to t in the auxiliary graph G_f that contains only edges with non-zero capacity.

Algorithm 1 FordFulkerson($G = (V, E, c)$)

1: Initialize $f(e) \leftarrow 0$ for all edges.
2: **while** $\exists$ augmenting path p in G_f **do**
3:     augment as much flow along p as possible.

Augmenting Path Algorithm

Animation for augmenting path algorithms is only available in the lecture version of the slides.

Augmenting Path Algorithm

Theorem 2

*A flow f is a maximum flow **iff** there are no augmenting paths.*

Theorem 3

The value of a maximum flow is equal to the value of a minimum cut.

Proof.

Let f be a flow. The following are equivalent:

1. There exists a cut A, B such that $\text{val}(f) = \text{cap}(A, B)$.

2. Flow f is a maximum flow.

3. There is no augmenting path w.r.t. f.

□

Augmenting Path Algorithm

1. $\implies$ 2.
This we already showed.

2. $\implies$ 3.
If there were an augmenting path, we could improve the flow.
Contradiction.

3. $\implies$ 1.

- ▶ Let f be a flow with no augmenting paths.
- ▶ Let A be the set of vertices reachable from s in the residual graph along non-zero capacity edges.
- ▶ Since there is no augmenting path we have $s \in A$ and $t \notin A$.

Augmenting Path Algorithm

$$\text{val}(f) = \sum_{e \in \text{out}(A)} f(e) - \sum_{e \in \text{into}(A)} f(e)$$

$$= \sum_{e \in \text{out}(A)} c(e)$$

$$= \text{cap}(A, V \setminus A)$$

This finishes the proof.

Here the first equality uses the flow value lemma, and the second exploits the fact that the flow along incoming edges must be 0 as the residual graph does not have edges leaving A.

Analysis

Assumption:
All capacities are integers between 1 and C.

Invariant:
Every flow value $f(e)$ and every residual capacity $c_f(e)$ remains integral troughout the algorithm.
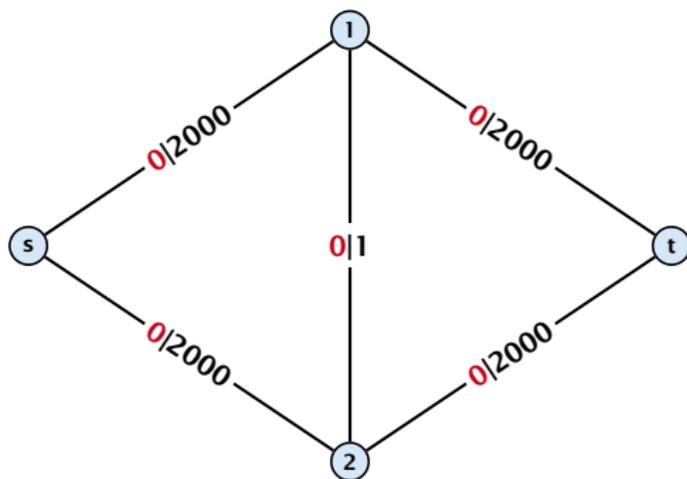
Lemma 4

The algorithm terminates in at most $\mathrm{val}(f^*) \le nC$ *iterations, where* f^* *denotes the maximum flow. Each iteration can be implemented in time* $\mathcal{O}(m)$. *This gives a total running time of* $\mathcal{O}(nmC)$.

Theorem 5

If all capacities are integers, then there exists a maximum flow for which every flow value $f(e)$ *is integral.*

A Bad Input

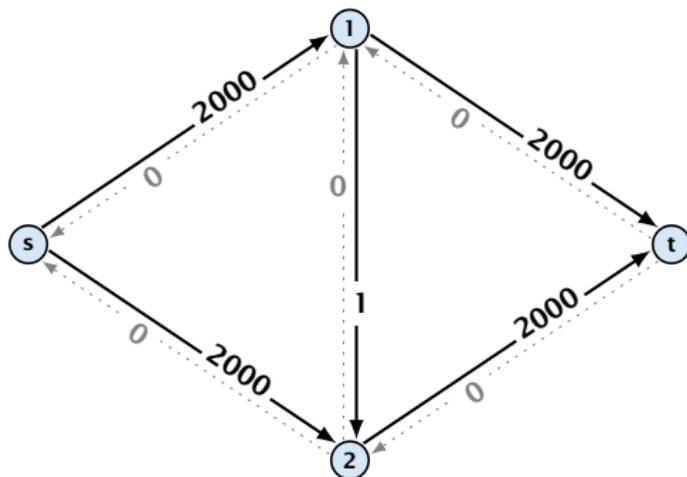Problem: The running time may not be polynomial.

Question:
Can we tweak the algorithm so that the running time is
polynomial in the input length?

A Bad Input

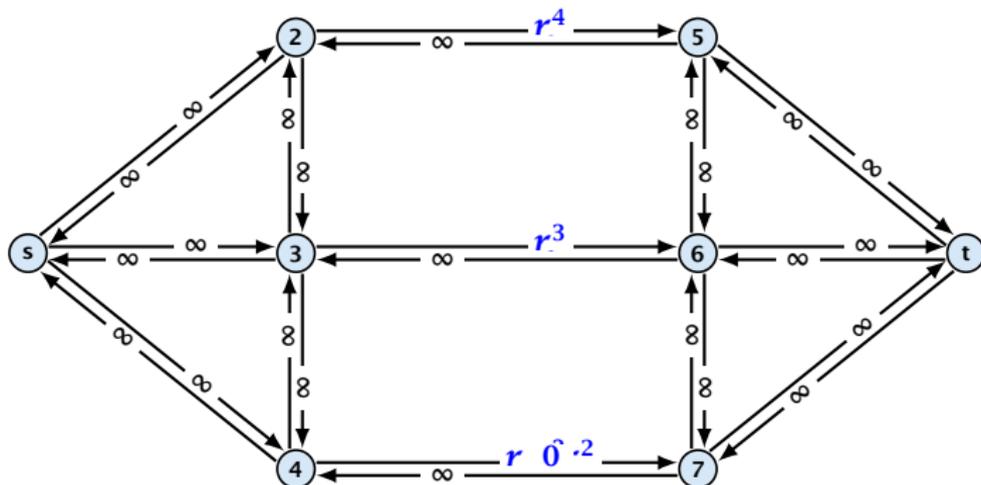Problem: The running time may not be polynomial.

Question:
Can we tweak the algorithm so that the running time is
polynomial in the input length?

See the lecture-version of the slides for
the animation.

A Pathological Input

Let $r = \frac{1}{2}(\sqrt{5} - 1)$. Then $r^{n+2} = r^n - r^{n+1}$.

Running time may be infinite!!!

See the lecture-version of the slides for the animation.

How to choose augmenting paths?

- ▶ We need to find paths efficiently.
- ▶ We want to guarantee a small number of iterations.

Several possibilities:

- ▶ Choose path with maximum bottleneck capacity.
- ▶ Choose path with sufficiently large bottleneck capacity.
- ▶ Choose the shortest augmenting path.